# Forking and categoricity in non-elementary model theory

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Let Elem(T), Emb(T), and Mod(T) be the respective categories. We will focus on Elem(T) at first, then will give a framework encompassing all of these and much more. Category-theoretic properties of Elem(T)

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- Non-examples: quantifier elimination (Morleyization gives an isomorphic category). Elimination of imaginaries (passing to T<sup>eq</sup> gives an equivalent category).
- Examples: categoricity in λ, stability in λ (but for this we need to describe cardinalities category-theoretically!).

# Presentability

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## Definition

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#### Definition (Gabriel-Ulmer, 1971)

An object M in a category  $\mathcal{K}$  is  $\lambda$ -presentable if for any  $\lambda$ -directed diagram  $D: I \to \mathcal{K}$ , any morphism  $f: M \to \operatorname{colim} D$  factors through some  $D_i$ :

$$M \longrightarrow \operatorname{colim} D$$

$$\bigcap_{D_i}$$

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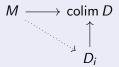
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The *(internal) size* of M is the predecessor of its presentability rank (if it exists).

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Shelah's thesis: dividing lines should have "internal" and "external" characterizations.

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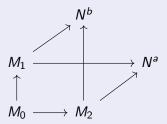
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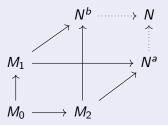
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Example: in  $Set_{mono}$ ,  $\{0\}$  and  $\{1\}$  have two non-equivalent amalgams over  $\emptyset$ :  $\{0, 1\}$  and  $\{1\}$  (with the obvious morphisms).

A stable independence notion is a class of squares (called independent squares, marked with  $\bot$ ) such that:

1. Independent squares are closed under equivalence of amalgam.

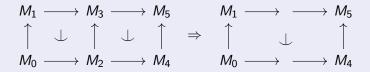
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- 4. Symmetry:

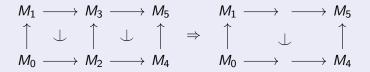
## Definition (stable independence notion - continued)

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 Accessibility: the category whose objects are arrows and whose morphisms are independent squares is accessible. This implies that any arrow can be "resolved" in an independent way:



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- In any accessible category with pushouts, the class of all squares form a stable independence notion.
- ► A first-order T is stable if and only if Elem(T) has a stable independence notion (given by nonforking squares).
- (LRV 2019) Assuming a large cardinal axiom, an accessible category with directed colimits and all morphisms monos has stable independence (on a cofinal full subcategory) if and only if it is stable in a proper class of cardinals.

# Stable independence and cofibrant generation

### Theorem (LRV)

Let  ${\cal K}$  be an accessible, bicomplete category. Let  ${\cal M}$  be a class of morphisms of  ${\cal K}$  such that:

- 1.  $\mathcal{M}$  contains all isomorphisms, is closed under transfinite compositions, pushouts, and retracts.
- 2. The induced category  $\mathcal{K}_{\mathcal{M}}$  is accessible and closed under directed colimits in  $\mathcal{K}$ .
- 3.  $\mathcal{M}$  is coherent: if  $M_0 \xrightarrow{f} M_1 \xrightarrow{g} M_2$ ,  $g, gf \in \mathcal{M}$ , then  $f \in \mathcal{M}$ .

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Then  $\mathcal{K}_{\mathcal{M}}$  has stable independence if and only if  $\mathcal{M}$  is cofibrantly generated (i.e. can be generated from a subset using transfinite compositions, pushouts, and retracts).

## New examples of stable independence

# Corollary (LRV)

- 1. The category of flat *R*-modules with flat morphisms (more generally, any AEC of "roots of Ext") has stable independence.
- 2. Any Grothendieck topos restricted to regular monos has stable independence.
- 3. Any Grothendieck abelian category restricted to regular monos has stable independence.
- Any combinatorial model category where all objects are cofibrant and whose cofibrations are coherent (e.g. monos) has stable independence, when restricted to its cofibrations.

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- ► Continuing in this fashion, we can inductively define when K has n-dimensional stable independence, for n ≥ 2.
- These *n*-dimensional diagrams are known in model theory (Shelah's main gap, Zilber's pseudoexponential fields, etc.) but the definitions are arguably complicated and ad-hoc.

#### Theorem (eventual categoricity)

Let  $\mathcal{K}$  be a finitely accessible category with all morphisms monos (or just an AEC). Assume either:

- 1. (Shelah-V.) There is a proper class of strongly compacts.
- 2. (V.)  $\Diamond_S$  for all stationary sets *S*, and *K* has no maximal objects.
- 3. (V. 2017)  $\mathcal{K}$  has "low complexity" (locally multipresentable/universal)

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### Open problems

- Can one prove eventual categoricity in ZFC, at least for finitely accessible categories or AECs?
- What is the status of eventual categoricity for accessible categories? Accessible categories with directed colimits?
- How do presentability ranks behave in general accessible categories? For example, can there be an object of limit presentability rank? The answer is no under SCH (LRV, 2019)
- When does existence of two-dimensional independence imply existence of higher dimensional independence?
- What are more examples of stable independence in mainstream mathematics (and computer science)?

# Thank you!

Some references:

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